

Descriptive Set Theory

Lecture 4

Prop. For a pruned tree T on A , $[T]$ is compact if and only if T is finitely-branching.

Proof \Rightarrow . Done.

\Leftarrow (Jenna Zouback). Let \mathcal{U} be an open cover of $[T]$ and suppose towards a contradiction that T is not finitely-branching. We call $s \in T$ **heavy** if $[T_s]$ cannot be covered by a finite subset of \mathcal{U} , where $T_s := \{t \in T : t \leq s \text{ or } t \geq s\}$. Thus, we know that \emptyset is heavy. Then one of the extensions of \emptyset must be heavy (finite union of finite sets is finite) and we continue, finding a branch $x \in [T]$ with the property that $[T_{x|_n}]$ cannot be covered by a finite subset of \mathcal{U} for each n . But \mathcal{U} covers $[T]$ so $\exists U \in \mathcal{U}$ with $x \in U$. But then for a large enough n , $[T_{x|_n}] \subseteq U$, a contradiction. \square

In particular, $2^{\mathbb{N}}$ is compact. Also, $\mathbb{N}^{\mathbb{N}}$ is not compact. In fact, $\mathbb{N}^{\mathbb{N}}$ is not even σ -compact (i.e. cbl union of compact).

unlike \mathbb{R}^n .

Prop. $\mathbb{N}^{\mathbb{N}}$ is not σ -compact.

Proof. Let $X \subseteq \mathbb{N}^{\mathbb{N}}$ be a σ -compact subset and we aim to find $y \in \mathbb{N}^{\mathbb{N}}$ that is not in X . Note that $X = \bigcup_{n \in \mathbb{N}} [T_n]$, where each T_n is a finitely-branching pruned tree.

For $x, y \in \mathbb{N}^{\mathbb{N}}$, we say that y **dominates** (resp. **eventually dominates**) x , if $\forall n \ x(n) \leq y(n)$ (resp. $\forall^\infty n, x(n) \leq y(n)$) ($\forall^\infty n$ means for all but finitely many, i.e. $\exists N \forall n \geq N$.)

($\exists^\infty n$ means for infinitely many, i.e. $\forall N \exists n \geq N$.)

By finite branching, each $[T_n]$ is dominated by an $x_n \in [T_n]$. By diagonalization, we get a $y \in \mathbb{N}^{\mathbb{N}}$ eventually dominating every x_n , namely, for each $i \in \mathbb{N}$,

let $y(i) := \max \{x_0(i), x_1(i), x_2(i), \dots, x_i(i)\}$. Indeed,

y eventually dominates every x_f because $\forall i \geq f$, $y(i) \geq x_f(i)$. Thus, $y \notin X$. □

The following dichotomy shows that $\mathbb{N}^{\mathbb{N}}$ is **the** non- σ -compact space:

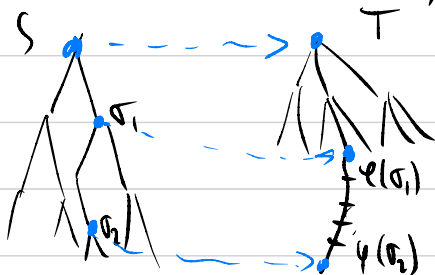
Hurewicz's dichotomy. Every Polish space X is either σ -compact or

X has a closed subset homeomorphic to $\mathbb{N}^{\mathbb{N}}$.



Monotone maps and continuous functions. We will learn how to build continuous functions $[S] \rightarrow [T]$, where S, T are trees, and we will see that all continuous functions are built this way.

Let S, T be trees on alphabets A, B , resp. Call a map $\varphi: S \rightarrow T$ monotone if for all $\sigma_1, \sigma_2 \in S$, $\sigma_1 \subseteq \sigma_2 \Rightarrow \varphi(\sigma_1) \subseteq \varphi(\sigma_2)$, and $\varphi(\emptyset) = \emptyset$.



Let $D_\varphi := \{x \in [S] : |\varphi(x|_n)| \rightarrow \infty \text{ as } n \rightarrow \infty\}$

Thus, φ induces a function

$$\begin{aligned} \varphi^*: D_\varphi &\rightarrow [T] \\ x &\mapsto \bigcup_n \varphi(x|_n) \end{aligned}$$

Theorem. (a) For any monotone $\varphi: S \rightarrow T$, D_φ is a G δ subset of $[S]$, and φ^* is continuous.
(b) All continuous maps from G δ subsets of $[S]$ to $[T]$ arise in this fashion.

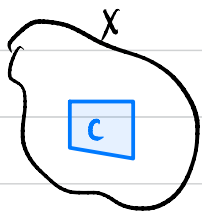
Proof. (b) is left for you.

(a) To see that D_φ is G δ , we note that $\forall x \in [S]$,
 $x \in D_\varphi \Leftrightarrow \forall \ell \in \mathbb{N} \exists n \in \mathbb{N} |\varphi(x|_n)| > \ell$.

$\underbrace{\qquad\qquad\qquad}_{\text{G}\delta} \underbrace{\qquad\qquad\qquad}_{\text{open}} \underbrace{\qquad\qquad\qquad}_{\text{closed}}$

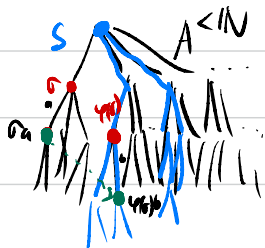
Note $\{x \in [S] : |\varphi(x|_n)| > l\}$ is a clopen set here membership in it depends only on first n coordinates. $\exists n \in \mathbb{N}$ is the same as $\bigcup_{n \in \mathbb{N}}$ and $\forall n \in \mathbb{N}$ is the same as $\bigcap_{n \in \mathbb{N}}$. \square

As a quick application, we consider retractions. A closed subset C of a topol. space X is called a **retract** of X if \exists continuous $\pi: X \rightarrow C$ s.t. $\pi|_C = \text{id}_C$. This map π is called a **retraction** of X onto C .



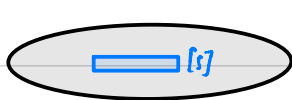
Cor. For any alphabet A , and closed subset $C \subseteq A^{\mathbb{N}}$ is a retract of $A^{\mathbb{N}}$. In particular, for closed sets $C_1 \subseteq C_2$, C_1 is a retract of C_2 .

Proof. Let $C = [S]$ for some pruned tree S on A .



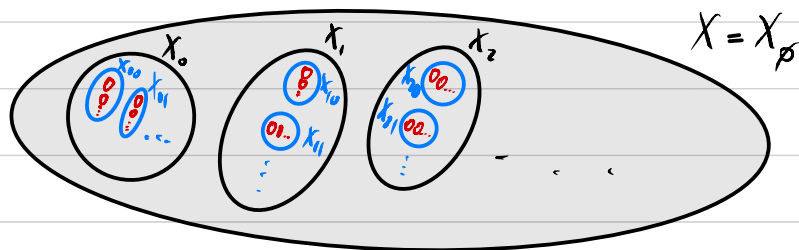
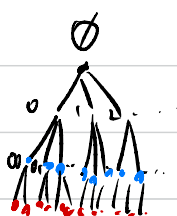
We define $\varphi: A^{<\mathbb{N}} \rightarrow S$ as follows:

$\varphi(\emptyset) := \emptyset$. Let $\sigma \in A^{<\mathbb{N}}$ s.t. $\varphi(\sigma)$ is defined but $\varphi(\sigma a)$ isn't defined yet. If $\sigma a \in S$, $\varphi(\sigma a) := \sigma a$. Otherwise, using that σ is pruned, take any $b \in A$ s.t. $\varphi(\sigma)b \in S$. Put $\varphi(\sigma a) := \varphi(\sigma)b$. \square



Luzin and Cantor schemes. A Luzin scheme in a top. space X is a sequence $(X_s)_{s \in \mathbb{N}^{<\mathbb{N}}}$ of subsets of X such that

- (i) $X_s \subseteq X_t$ if $t \geq s$, $\forall s, t \in \mathbb{N}^{<\mathbb{N}}$.
- (ii) $X_{sa} \cap X_{sb} = \emptyset$ $\forall s \in \mathbb{N}^{<\mathbb{N}}$ and $a \neq b \in \mathbb{N}$.



- (iii) If X is a metric space with a metric d , then we'd say that the scheme has vanishing diameter if $\forall x \in \mathbb{N}^{<\mathbb{N}}$, $\text{diam}_d(X_{x|n}) \rightarrow 0$ as $n \rightarrow \infty$.

A scheme with vanishing diameter induces a function $f: D \rightarrow X$, where $D := \{x \in \mathbb{N}^{<\mathbb{N}} : \bigcap_n X_{x|n} \neq \emptyset\}$.
 $x \mapsto$ the unique element of $\bigcap_n X_{x|n}$ (by vanishing diam).

If instead we have $(X_s)_{s \in \mathbb{Z}^{<\mathbb{N}}}$, then we call it a Cantor scheme.

Properties of Luzin schemes. let $(A_s)_{s \in \mathbb{N}^{<\mathbb{N}}}$ a Luzin scheme of vanishing diameter in a metric space (X, d) ,
and let $f: \mathcal{D} \rightarrow X$ be the induced map.

(a) f is injective and continuous.

(b) If $A_s = \bigcup_{n \in \mathbb{N}} A_{s \smallfrown n} \quad \forall s \in \mathbb{N}^{<\mathbb{N}}$, then f is surjective.

(c) If A_s is open for each $s \in \mathbb{N}^{<\mathbb{N}}$, then f is open.

(d) If d is complete and $\overline{A_{s \smallfrown n}} \subseteq A_s \quad \forall s \in \mathbb{N}^{<\mathbb{N}}$ and $n \in \mathbb{N}$, then \mathcal{D} is closed. In fact, $x \notin \mathcal{D} \iff \exists n \quad A_{x \smallfrown n} = \emptyset$.

In particular, if all $A_s \neq \emptyset$, then $\mathcal{D} = \mathbb{N}^{\mathbb{N}}$.